

Tutorial session 2: Programming with Maxima
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Recall:

- (a) The **for** loop provides the ability to execute a statement repeatedly for a fixed number of times.
- (b) If we want to repeat a statement, but we do not know how many times, we can use the **while** loop. The expression between **while** and **do** has to be a boolean expression (an expression returning **true** or **false**).
- (c) Sometimes we want to make a case-by-case analysis. For this purpose we can use the **if** statement. The statements after **then** are only evaluated, if the condition is true.

1. Define the list $L = [1, 2, \dots, 1000]$.
 - (a) Derive the list $L1 = [2^p + 1, p \in L]$.
 - (b) Count the number of primes in $L1$.
 - (c) Derive the list of all the prime elements in $L1$.
 - (d) Derive the list $L2 = [p \in L / 2^p + 1 \text{ is prime}]$.
 - (e) Check that each element in $L2$ can be written as power of 2.
 - (f) Compute the arithmetic mean of the elements of L and return its numerical approximation.
2. Let $a, b \in \mathbb{N}_{\geq 0}$. The division with remainder gives the relation

$$a = bq + r, \quad 0 \leq r < b. \quad (1)$$

Use the fact that $\gcd(a, b) = \gcd(b, r)$ and $\gcd(a, 0) = a$ to implement the function `gcd1(a, b)` which computes recursively the greatest common divisor of a and b . Hint: The remainder r in (1) is `mod(a, b)` in Maxima.

Test your program for the computation of $\gcd(2^{400} + 3; 3^{300} + 8)$ and compare the timings with the internal function `gcd`.

3. Consider the matrix

$$M = \begin{pmatrix} 1 & 2 & 3 \\ 4 & 5 & 5 \\ 7 & 8 & 9 \end{pmatrix}.$$

- (a) Write a first recursive program `MatPow1(A, n)` to compute A^n ($n = 0, 1, 2, \dots$) using the relation $A^n = A \times A^{n-1}$. Apply this with M^{10} and M^{1000} (use `.` for matrix multiplication and `ident(n)` to get the identity matrix of order n).
- (b) Write the second recursive program `MatPow2(A, n)` to compute A^n ($n = 0, 1, 2, \dots$) using the divide-and-conquer approach: $A^n = (A^2)^{n/2}$ if n is even and $A^n = A \times A^{n-1}$ if n is odd.

For each of the implementations, compare the timings with the internal function `A^^n` for $n = 10000$ for example.

4. Implement the procedures `PolyQuot(a, b, x)` and `PolyRem(a, b, x)` which compute, respectively, the quotient and the remainder of $a(x)$ by $b(x)$ for two polynomials $a(x), b(x)$. Use your function to find the polynomial quotient and remainder of the division of $a(x) = 12x^6 - 8x^5 + 17x^4 + 5x^3 - 4x^2 + 3x + 5$ by $b(x) = 3x^4 + 5x - 1$ and $b(x)$ by $a(x)$ and check your results with the internal command `divide(a, b)`.

Hint: `coeff(p, x, n)` returns the coefficient of x^n in the polynomial p , `hipow(p, x)` returns the degree of the polynomial p .

5. The Chebyshev polynomials $T_n(x)$ are defined by

$$T_n(x) = 2xT_{n-1}(x) - T_{n-2}(x), \quad T_0(x) = 1 \text{ and } T_1(x) = x; \quad (2)$$

and have the property

$$\begin{cases} T_n(x) = 2(T_{\frac{n}{2}}(x))^2 - 1, & \text{if } n \text{ is even} \\ T_n(x) = 2T_{\frac{n-1}{2}}(x)T_{\frac{n+1}{2}}(x) - x, & \text{if } n \text{ is odd} \\ T_0(x) = 1, \quad T_1(x) = x. \end{cases} \quad (3)$$

Implement each of the above recurrence equations and compare their complexities by taking higher values of n .